

View-Frustum Culling



 In many real scenes, a substantial percentage of the scene is outside the view frustum

> Potentially Visible Set







Bounding Volumes (BVs)



- Test per polygon is too expensive, overall rendering time would be slower than without VFC
- Therefore, test complete objects (= set of polygons) whether they are outside the view frustum
- Do fast tests with simple bounding volumes (BVs):





Axis Aligned Bounding Box (AABB)



Oriented BBox (OBB)

- The process is efficient only if
 - Cost(BV test) << Cost(rendering the polygon set)



Calculation of OBBs









AABB := (min, max) =
 (x_{min}, y_{min}, z_{min}, x_{max}, y_{max}, z_{max})



- OBB is defined by
 - center
 - 3 axes
 - 3 "radii"
 - Corresponds to a 3x4 matrix:
 - $T(M) \cdot R(u,v,w) \cdot S(r_x,r_y,r_z)$







Representation of the View Frustum

- Procedure:
 - 1. Use parameters from gluPerspective and gluLookAt
 - 2. Calculate vertices of the frustum
 - 3. Calculate the frustum planes
- Determine corners (in world coordinates):

$$F = C + f \cdot \mathbf{d}$$

 $P = F + \frac{1}{2}H\mathbf{v} - \frac{1}{2}W\mathbf{u}$

Analogously, calc all other vertices

- Determine the planes of the corners :
 - 3 points are sufficient (cross product of edges)
 - Note: ensure a consistent orientation of the normals!
 - Small optimization: normals of the near and far plane are known already





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Given: 6 plane equations

$$E_i:x\cdot n_i-d_i=0$$

and a sphere

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$$\left(x-c\right)^2-r^2=0$$

Question: Is the sphere is completely outside the frustum?

• Yes
$$\leftrightarrow \exists i : c \cdot n_i - d_i > r$$

If
$$\exists i: |c \cdot n_i - d_i| \leq r$$

then one of the planes intersects the sphere (but not necessarily the frustum)

• If
$$\forall i: c \cdot n_i - d_i < -r$$

then the sphere is completely inside the frustum





Test Box v. Frustum



- Warning: it is not sufficient to check that all vertices are outside the frustum!
 - Counterexample:
- A simple, conservative test:
 All 8 vertices are on the positive side of the same plane → box is outside
- This test produces so-called
 "false positives" → increases the PVS
- The box is completely inside ⇔
 all vertices are on the negative side of all planes







Optimizations

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 We denote by "N vertex" that vertex of all vertices where the following function assumes the minimum:

 $f(\mathbf{x}) = \mathbf{x} \cdot \mathbf{n} - \mathbf{d}$

- Analogously define "P vertex" (f assumes max)
- These are (almost always) unique because f is monotone, and a box is convex







- How to quickly find the N- or P vertex?
- If box = axis-aligned bounding box (AABB), then it can be done very fast
- $AABB = (x_{min}, y_{min}, z_{min}, x_{max}, y_{max}, z_{max})$





Further Optimizations



 "Meta-BVs": If many boxes need to be tested, enclose boxes and balls in a frustum



• Or enclose the frustum in an AABB, too



- Produces more false positives, so YMMV
- Exploit temporal coherence: if box has been culled by a certain frustum plane, save that plane and test this *first* in the next frame.
 Probability is high that this plane culls the box again!



Hierarchical View Frustum Culling



■ Generating at each node of the scene graph a bounding volume including the complete subtree →

Bounding-Volume-Hierarchy (BVH)

Traverse this BVH und test all knots





Further Optimization

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Plane Masking:

- If a box is completely on the negative side of a plane, then all children too. Do not test this level for the children
- If a BV is completely inside, then all children are inside









- Occlusion Culling is always interesting, if many objects are hidden by a few objects
- Definition: Depth Complexity
 - Number of intersections of the ray in the scene
 - Number of polygons projected on a pixel
 - Number of polygons that would be visible at a pixel, all polygons were transparent
- Comment : Depth Complexity is oberservation and directional





Examples of High Depth Complexity







First, the Special Case of "Cities"



- Render the scene from front to back (reverse Painter's Algorithm)
- Generate an "occlusion horizon"







- Rendering an object (here tetrahedra; behind the gray objects):
 - Determine axis-aligned bounding box (AABB) of the projection of the object
 - Comparison with the occlusion horizon







- If an object is considered as visible:
 - Add the AABB with the previous occlusion horizon



General Occlusion Culling



Given:

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- A partial(!) rendered scene, and
- not yet rendered object
- Task:
 - Decide quickly whether the object would modify pixels in the frame buffer, if it were rendered;
 - I.o.w.: decide quickly whether the object of the current scene is completely covered
- Terminology:
 Occluder
 Occluded geometry ("occludee")



Examples of Applications of the General Occlusion Culling

















Invisible polygons: 10M (ca. 96%)



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- Earlier as extension ARB_occlusion_query , nowadays in OpenGL core from version 1.5
 - Operating mode: Asks OpenGL how many pixels would be "repainted" from a batch
- Appendage: Draw a simple representation ("Proxy"), without changing the color or depth buffer
 - Were no pixels drawn by the proxy, the object itself must not be drawn
- Proxy geometry: first sacrifice a little computing power to possibly save a lot of computing power later
 - Tolerably accurate bounding volumes
 - No texturing, no shading, no light sources
 - No colors, texture coordinates, normals





• First create occlusion query at initialization :

glGenQueries(int count, unsigned int queryIDs[]);

- Render a set of objects (hiding a lot)
- Disable writing in Z- and color buffer (optional):

```
glDepthMask( GL_FALSE );
glColorMask( GL FALSE,GL FALSE,GL FALSE,GL FALSE );
```

Create request for a lot of other objects :

```
glBeginQuery( GL_SAMPLES_PASSED, unsigned int querynum );
// rendere Proxy-Geometrie, z.B. Bounding Volume ...
glEndQuery( GL_SAMPLES_PASSED );
```

Reading result of the request:



Demo



0	e e occlusion_query.cpp (~/Work/Lehre/CG1/demos/occlusion_query) - VIM
voi	d draw_objects()
{	glColor3f(1,1,0); glPushMatrix(); glTranslatef(0,025, 0); glScalef(1, .05, 1);
	<pre>// render cube, with occlusion query glBeginQueryARB(GL_SAMPLES_PASSED_ARB, oq_plane); glutSolidCube(.5); glEndQueryARB(GL_SAMPLES_PASSED_ARB); glPopMatrix();</pre>
}	<pre>// render sphere, with occlusion query glColor3f(1, 0, 0); glPushMatrix(); glTushMatrix(); glBeginQueryARB(GL_SAMPLES_PASSED_ARB, oq_sphere); glutSolidSphere(.25, 20, 20); glEndQueryARB(GL_SAMPLES_PASSED_ARB); glPopMatrix();</pre>
voi	<pre>set_app_info_string()</pre>
{	GLuint plane_samples, sphere_samples;
	<pre>// get results of occlusion queries glGetQueryObjectuivARB(oq_plane, GL_QUERY_RESULT_ARB, &plane_samples); glGetQueryObjectuivARB(oq_sphere, GL_QUERY_RESULT_ARB, &sphere_samples);</pre>
	<pre>string s; char buff[80];</pre>
	<pre>s = "visible samples\n plane: "; sprintf(buff, "%d", plane_samples); s += buff; if(plane_samples == 0) { s += " no samples visible"; } s += "\n sphere: ";</pre>
	<pre>sprintf(buff, "%d", sphere_samples); s += buff;</pre>
	T(sphere_samples == 0)
	s += " no samples visible";



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- Problem: A query = expensive State-Changes
 - Before: Disable writing to color-and Z-buffer
 - After: Switch back
 - This overhead takes more time than the actual query!
- Idea: Batching
- Implement 2 additional queues
 - Both contain objects that should to be tested for visibility
 - I-Queue: contains previously "invisible" objects
 - V-Queue: likewise for "visible"
 - Parameters: Batch size b (ca. 20-80)
 - Principle: only if batch size is reached, the list of queries is sent to OpenGL
- Previously visible objects are still rendered immediately





Exemple: each color = one state change



Naive

CHC++





- Goal: Reduce the number state changes, and thus the time required per Occlusion Query
- Therefore, send a sequence of requests, read the result of the sequence afterwards





The Naive "draw-and-wait" Approach



Sort items about the depth in Create query sequence while some objects are not rendered: For each object in query sequence: BeginQuery Render bounding volume EndQuery For each object in query sequence: GetQuery if #pixel drawn > 0: Render object





- Problems of the naive approach:
 - Very high response time (latency) for a query:
 - long graphics pipeline,
 - some time by the execution of the queries (rasterization), and
 - transfer the result back to the host.



Sequence: "CPU stalls" and "GPU starvation"



Sort the Object List



 Observation: Depending on the order in which you render the objects, you get a high culling rate or not



Solution: sort the list by distance to the object Viewpoint







• Often only conservative culling :

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- Even if only one pixel of the BVs is visible also one pixel of the object can be visible → draw object
- Disadvantage: Often outer parts of the BVs are visible where no object pixel are located
- Idea: Ignore barely visible objects
 - Object probably (!) not visible if only a few pixels of the BVs are visible
 - Heuristics: Draw object only if query result ≥ threshold
 - Potentially "small" holes in or between objects









- Here in a simplified representation (a.o. without hierarchy)
- Given: Set of objects
 - Here: Object = Amount of useful contiguous polygons
- Ideas:

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- Perform a queue with stored hardware occlusion queries
- First assumption: if an object was visible in the last frame, it is also visible in the current frame
- If an object was invisible, first check its visibility
- Do not wait for result, go further through the query
- Edit query results as soon as they are available





```
L = list of all objects (incl. BVs)
Q = queue for occlusion queries (initally empty)
sort L from front to back with respect to current viewpoint
repeat:
  // process list of objects
  if L not empty:
    O = L. front
    if O inside view frustum:
      issue occlusion query with BV(O)
      append 0 to Q
      if O is marked "previously visible":
        render O
  end if
. . .
```





```
. . .
  // process queries
  while Q not empty and
        result of occlusion query Q.front available
    V = Q.pop
    if num. visible pixels of query V > threshold:
      V.obj = "visible"
      if V.obj is not marked "previously visible":
          render V.obj
    else:
      V.obj = "invisible"
  end while
until Q empty and L empty
```

Below: gradual improvement of this algorithm

Fusion (Potentially) Hidden Geometry



Observation :

- If we knew that a lot of objects in the current frame is hidden, then we could verify this by exactly one occlusion query
- Objects that were hidden in many frames are probably obscured in the current frame (*temporal coherence of visibility*)
- Idea:

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- Invent an "oracle" that can predict for a given set of objects with high probability whether the coherence of visibility is satisfied
- If the probability is high enough, test this set by 1

Query:

```
glBeginQuery( GL_SAMPLES_PASSED, q );
render BVs with the set of objects ...
glEndQuery( GL_SAMPLES_PASSED );
glGetQueryObjectiv( q, GL_QUERY_RESULT, *samples );
```

This will be called in the following: Multiquery!





Definition: Visibility persistence

$$p(t) = \frac{I(t+1)}{I(t)}$$

where I(t) = number of objects, which were constantly covered in previous t frames

- Interpretation: p(t) = " probability ", that one object, which was covered t frames, will also be covered in the following frames
- Observation: is amazingly
 independent from object and scene
 Observation: is amazingly
- Consequence: can be approximated well by analytic function!

$$p(t)\approx 0.99-0.7e^{-t}$$





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- If t_O = Number of previous frames which the object O was covered
- Define an "oracle" for a set M of objects *i(M)* := the "probability" that all objects from M in the actual frame will be invisible (only a heuristic!):

$$i(M) = \prod_{O \in M} p(t_O)$$

- Define that
 - Costs of an occlusion multiquery (in the batch):

$$C(M) = 1 + c_1 |M|$$

• Expected benefits of a multiquery: $B(M) = c_2 i(M) \sum_{O \in M} \text{num polygons of } O$





Thus defining the expected value of a multiquery:

$$V(M) = \frac{B(M)}{C(M)}$$

- f the I-queue is full at some point:
 - Sort the objects O_i in the I-queue $t_O \rightarrow \{O_1, \dots, O_b\}$
 - Simple greedy search the maximum

$$\max_{n=1\dots b} \left\{ V(\{O_1,\dots,O_n\} \right\}$$

- Set on a multiquery for these first *n* objects from the I-queue
- Repeat until the I-queue is empty





- Observation: As greater the BV in relation to the object as more probable that the occlusion query returns a "*false positive*" (claims "visible", but in truth "invisible")
- Objective: close as possible BVs
- Boundary conditions:
 - BVs must be very fast to render
 - BVs may not cost a lot of memory
- Idea:

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- Decompose object into cluster (cluster of polygons)
- Wrap a BBox around each clustert (AABB)
- Use as BV foreach object the union of the small BBoxes







- Question: how small should the "small" AABBs (or cluster) be?
- Observation: the greater the number of small AABBs, ...
 - ... the greater the probability that "invisible" is correctly recognized, but
 - ... the greater the surface → onger rendering time of the resulting occlusion queries
- Strategy for the construction of the "narrow AABBs":
 - Devide clusters recursively
 - termination criterion: if \sum Oberfläche der kleinen AABBs > $\sigma \cdot$ Oberfläche der großen AABB
 - Parameter σ depends on the graphics card ($\sigma \approx 1.4$ seems OK)



Altogether



• The queues in CHC++:







Walkthrough the power plant model:







State Changes: CHC vs. CHC++

Each color represents a state change required by the algorithm





Powerplant Walkthrough 2





CHC++ Multiqueries

Each color represents nodes covered by a single multiquery





Coherent Hierarchical Culling Hardware Occlusion Queries Made Useful

Culling 76



Another Special Case: Architectural Models





Cells and Portals (Portal Culling)



- Scenario: Walkthrough of buildings and cities
- Transparent portals connect the cells
 - doors, windows, holes, ...

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Observation: cells see each other only through the portals



- Which cell is included in the PVS?
 - The cell which contains the Viewpoint
 - And these cells, which have a portal to the initial cell



Example scene











Example scene :



- Speedup is highly dependent on the model and viewpoint
 - Frame rate is 1-10-fold the frame rate without Cells & Portals method
 - For typical viewports the method removes 20 50% from the model



- Field of applications
 - Computer games
 - Buildings
 - Cities
 - Ships (inside)
- Not suitable for CAD data
 - Aircrafts
 - Industrial facilities
- Not suitable for natural objects
 - Plants
 - Forests







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- Idea: Objects that occupy less than N pixels in the projection are not shown
- This approach also removes parts that would be visible in the final image
- Advantage: trade-off quality / speed





Particularly suitable for camera motion (the faster, the more details can be culled)



• Estimate the size of the BVs in screen space from:

